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Extending the CLP engine for reasoning under uncertainty

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talk structure

- Logic Programming (LP)
- Uncertainty and LP
- Constraint LP
- clp(pfd(Y))
- clp(pfd(c))
- Caesar's coding experiments

Used in AI for crisp problem solving and for building executable models and intelligent systems.

Programs are formed from logic based rules.

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member( H, [H|T] ).
member( EI, [H|T] ) :- member( EI, T ).
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execution tree



member(H, [H|T]). member(EI, [H|T]) :- member(EI, T).

uncertainty in logic programming

Most approaches use Probability Theory but there are fundamental questions unresolved.

In general,

0.5 : member(H, [H|T]). 0.5 : member(El, [H|T]) :- member(El, T).

stochastic tree



0.5 : member(H, [H|T]). 0.5 : member(EI, [H|T]) :- member(EI, T).

constraints in lp

Logic Programming :

- execution model is inflexible, and
- its relational nature discourages use of state information.

constraints in lp

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Constraints add

- specialised algorithms

constraints in lp

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Constraints add

- specialised algorithms
- state information

constraint store

?- Q. X # Y

Logic Programming engine

Constraint store interaction

constraints inference



finite domain distributions

For discrete probabilistic models clp(pfd(Y)) extends the idea of finite domains to admit distributions.

from clp(fd)

 $X in \{a, b\}$ (i.e. X = a or X = b)

to clp(pfd(Y))

$$p(X = a) + p(X = b)$$

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from clp(fd)

 $X in \{a, b\}$ (i.e. X = a or X = b)

to clp(pfd(Y))

$$[p(X = a) + p(X = b)] = 1$$

constraint based integration

Execution, assembles the probabilistic model in the store according to program and query.

Dedicated algorithms can be used for probabilistic inference on the model present in the store.

probability of predicates

- pvars(E) variables in predicate E,
- e vector of finite domain elements
- $p(e_i)$ probability of element e_i
- S a constraint store.
- E/e E with variables replaced by e.

The probability of predicate E with respect to store S is

$$P_{\mathcal{S}}(E) = \sum_{\substack{\forall e \\ \mathcal{S} \vdash E/e}} P_{\mathcal{S}}(e) = \sum_{\substack{\forall e \\ \mathcal{S} \vdash E/e}} \prod_{i} p(e_i)$$



Probabilistic variable definitions

 $Coin \in_{\mathbf{p}} finite_geometric([head, tail], 2)$

Conditional constraint

 $Dependent ~ {\tt I} {\tt \ddagger} ~ Qualifier$



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In clp(fd) systems labelling instantiates a group of variables to elements from their domains.

In clp(pfd(Y)) the probabilities can be used to guide labelling. For instance we have implemented

label(Vars, mua, Prbs, Total)

Selector *mua* approximates best-first algorithm which instantiates a group of variables to most likely combinations.



Each letter is encrypted to a random letter. Words drawn from a dictionary are encrypted. Programs try to decode them. We compared a clp(fd) solution to clp(pfd(c)).

clp(fd) no probabilistic information, labelling in lexicographical order.

clp(pfd(c)) distributions based on frequencies, labelling using mua.

proximity functions

- X_i - variable for ith encoded letter

- D_i dictionary letter
- freq() frequency of letter

$$px(X_i, D_j) = \frac{1/|freq(X_i) - freq(D_j)|}{\sum_k 1/|freq(X_i) - freq(D_k)|}$$

execution times



clp(pfd(c)) and clp(fd) on SICStus 3.8.6



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Constraint LP based techniques can be used for frameworks that support probabilistic problem solving.

clp(pfd(Y)) can be used to take advantage of probabilistic information at an abstract level.



Three curtains hiding a car and two goats. Contestant chooses an initial curtain. A close curtain opens to reveal a goat. Contestant is asked for their final choice.

What is the best strategy ? Stay or Switch ?

Monty Hall solution

If probability of switching is Swt, (Swt = 0 for strategy Stay and Swt = 1 for Switch) then probability of win is $P(\gamma) = \frac{1+Swt}{3}$.

Monty Hall in clp(cfd(c))

curtains(gamma, Swt, Prb) :- $Gift \in_{\mathbf{p}} uniform([a, b, c]),$ $First \in_{\mathbf{p}} uniform([a, b, c]),$ $Reveal \in_{\mathbf{p}} uniform([a, b, c]),$ $Second \in_{\mathbf{p}} uniform([a, b, c]),$ $Reveal \neq Gift, Reveal \neq First, Second \neq Reveal,$ $Second \ \mathbf{ls}_{wt} \ First,$ $Prb \ is \ \mathbf{p}(Second=Gift).$

Strategy γ Query

Querying this program ?- curtains(gamma, 1/2, Prb) Prb = 1/2.

Strategy γ Query

Querying this program ?- curtains(gamma, 1/2, Prb) Prb = 1/2.

?- curtains(gamma, 1, Prb) Prb = 2/3.

Strategy γ Query

Querying this program ?- curtains(gamma, 1/2, Prb) Prb = 1/2.

?- curtains(gamma, 1, Prb) Prb = 2/3.

?- curtains(gamma, 0, Prb) Prb = 1/3.



Other discrete probabilistic inference is possible. For instance Bayesian Networks representation and inference.



A B C

	A = y	A = n
C = y	0.60	0.90
C = n	0.40	0.10

clp(pfd(BN)) program





Zy = 0.66

current inference scheme



bottom line

