

Example of Constructing a Predictive Parsing Table

Grammar:

$$\begin{aligned} \underline{e} &\rightarrow \underline{t} \underline{e}' \\ \underline{e}' &\rightarrow + \underline{t} \underline{e}' \\ &| \varepsilon \\ \underline{t} &\rightarrow \underline{f} \underline{t}' \\ \underline{t}' &\rightarrow * \underline{f} \underline{t}' \\ &| \varepsilon \\ \underline{f} &\rightarrow (\underline{e}) \\ &| x \\ &| y \end{aligned}$$

first and ***follow*** sets:

Non-Terminal	<i>first</i>	<i>follow</i>
\underline{e}	'(', 'x', 'y'	\$, ')'
\underline{e}'	+', ε	\$, ')'
\underline{t}	'(', 'x', 'y'	+', \$, ')'
\underline{t}'	*, ε	+', \$, ')'
\underline{f}	'(', 'x', 'y'	*, '+, ')', \$

Start with an empty parsing table; the **rows** are non-terminals and the **columns** are terminals.

Non-Terminal	Input Symbol						
	x	y	$+$	$*$	$($	$)$	$\$$
\underline{e}							
\underline{e}'							
\underline{t}							
\underline{t}'							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{e} \rightarrow \underline{t} \underline{e}'$

$\text{first}(\underline{t} \underline{e}') = '(', 'x', 'y'$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}					$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'							
\underline{t}							
\underline{t}'							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\epsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{e} \rightarrow \underline{t} \underline{e}'$

$\text{first}(\underline{t} \underline{e}') = '(, 'x', 'y'$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$					$\underline{e} \rightarrow \underline{t} \underline{e}'$	
\underline{e}'							
\underline{t}							
\underline{t}'							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\epsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{e} \rightarrow \underline{t} \underline{e}'$

$\text{first}(\underline{t} \underline{e}') = \text{'('}, \text{'x'}, \text{'y'}$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$				$\underline{e} \rightarrow \underline{t} \underline{e}'$	
\underline{e}'							
\underline{t}							
\underline{t}'							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{e}' \rightarrow + \underline{t} \underline{e}'$

$\text{first}(+ \underline{t} \underline{e}') = '+'$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$				$\underline{e} \rightarrow \underline{t} \underline{e}'$	
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$				
\underline{t}							
\underline{t}'							
f							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{e'} \rightarrow \varepsilon$

$\text{follow}(\underline{e'}) = \{ \$, '}'$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e'}$	$\underline{e} \rightarrow \underline{t} \underline{e'}$				$\underline{e} \rightarrow \underline{t} \underline{e'}$	
$\underline{e'}$			$\underline{e'} \rightarrow + \underline{t} \underline{e'}$				$\underline{e'} \rightarrow \varepsilon$
\underline{t}							
$\underline{t'}$							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{e}' \rightarrow \varepsilon$

$\text{follow}(\underline{e}') = \$, \text{'}'$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$				$\underline{e} \rightarrow \underline{t} \underline{e}'$	
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}							
\underline{t}'							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{t} \rightarrow \underline{f} \underline{t}'$

$\text{first}(\underline{f} \underline{t}') = \text{'('}, \text{'x'}, \text{'y'}$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}					$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{t} \rightarrow \underline{f} \underline{t}'$

$\text{first}(\underline{f} \underline{t}') = '(', 'x', 'y'$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}	$\underline{t} \rightarrow \underline{f} \underline{t}'$				$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
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Current production:

$\underline{t} \rightarrow \underline{f} \underline{t}'$

$\text{first}(\underline{f} \underline{t}') = \text{'('}, \text{'x'}, \text{'y'}$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}	$\underline{t} \rightarrow \underline{f} \underline{t}'$	$\underline{t} \rightarrow \underline{f} \underline{t}'$			$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'							
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{t'} \rightarrow \alpha$

$\text{follow}(\underline{t'}) = \{ '+', \$, ')' \}$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}	$\underline{t} \rightarrow \underline{f} \underline{t}'$	$\underline{t} \rightarrow \underline{f} \underline{t}'$			$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'			$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow * \underline{f} \underline{t}'$			
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{t'} \rightarrow \alpha$

$\text{follow}(\underline{t'}) = \text{'+'}, \text{'$'}, \text{'('}$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}	$\underline{t} \rightarrow \underline{f} \underline{t}'$	$\underline{t} \rightarrow \underline{f} \underline{t}'$			$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'			$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow * \underline{f} \underline{t}'$			$\underline{t}' \rightarrow \varepsilon$
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{t'} \rightarrow \alpha$

$\text{follow}(\underline{t'}) = \text{'+'}, \text{'$'}, \text{'('}$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}	$\underline{t} \rightarrow \underline{f} \underline{t}'$	$\underline{t} \rightarrow \underline{f} \underline{t}'$			$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'			$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow * \underline{f} \underline{t}'$		$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow \varepsilon$
\underline{f}							

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$f \rightarrow (\underline{e})$

$\text{first}(\underline{e}) = '('$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}	$\underline{t} \rightarrow \underline{f} \underline{t}'$	$\underline{t} \rightarrow \underline{f} \underline{t}'$			$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'			$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow * \underline{f} \underline{t}'$		$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow \varepsilon$
\underline{f}					$\underline{f} \rightarrow (\underline{e})$		

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{f} \rightarrow x$

$\text{first}(x) = \underline{'x'}$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}	$\underline{t} \rightarrow \underline{f} \underline{t}'$	$\underline{t} \rightarrow \underline{f} \underline{t}'$			$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'			$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow * \underline{f} \underline{t}'$		$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow \varepsilon$
\underline{f}	$\underline{f} \rightarrow x$				$\underline{f} \rightarrow (\underline{e})$		

for each production $\underline{n} \rightarrow \alpha$
for each $a \in \text{first}(\alpha)$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$
if $\varepsilon \in \text{first}(\alpha)$ then
for each $b \in \text{follow}(\underline{n})$
 add $\underline{n} \rightarrow \alpha$ to $T[\underline{n}, a]$

Current production:

$\underline{f} \rightarrow y$

$\text{first}(y) = \underline{'y'}$

Non-Terminal	Input Symbol						
	x	y	+	*	()	\$
\underline{e}	$\underline{e} \rightarrow \underline{t} \underline{e}'$	$\underline{e} \rightarrow \underline{t} \underline{e}'$			$\underline{e} \rightarrow \underline{t} \underline{e}'$		
\underline{e}'			$\underline{e}' \rightarrow + \underline{t} \underline{e}'$			$\underline{e}' \rightarrow \varepsilon$	$\underline{e}' \rightarrow \varepsilon$
\underline{t}	$\underline{t} \rightarrow \underline{f} \underline{t}'$	$\underline{t} \rightarrow \underline{f} \underline{t}'$			$\underline{t} \rightarrow \underline{f} \underline{t}'$		
\underline{t}'			$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow * \underline{f} \underline{t}'$		$\underline{t}' \rightarrow \varepsilon$	$\underline{t}' \rightarrow \varepsilon$
\underline{f}	$\underline{f} \rightarrow x$	$\underline{f} \rightarrow y$			$\underline{f} \rightarrow (\underline{e})$		