Using inference to improve search on SAT problems

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Introduction

- SAT: propositional satisfiability problem
- Archetypal NP-complete problem (Cook, 1971)
- Numerous practical applications, including planning, quasigroup completion, and model checking
- My research: evaluating techniques for using inference to improve search

Satisfiable SAT instance

$$\begin{array}{cccc} a & \vee & \neg b \\ b & \vee & c \\ \neg c & \vee & d \end{array}$$

Alternative form:

$$\Sigma = (a \vee \neg b) \wedge (b \vee c) \wedge (\neg c \vee d)$$

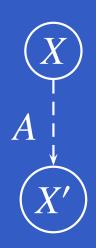
Satisfying assignment:

$$A = \{a \mapsto T, b \mapsto T, c \mapsto F\}$$

Solving SAT problems

- The Davis-Putnam method
 - Uses resolution to eliminate variables
- Resolution is intractable and impractical for many problems
 - E.g. *best* case exponential space complexity on pigeon-hole problems
- DLL (Davis-Logemann-Loveland)
 - Replaces resolution with branching
 - Complete backtracking search
 - Basis of the fastest available solvers

Assignment during search



$$X = (a \lor \neg b \lor \neg c)$$

$$A = \{a \mapsto F, b \mapsto T\}$$

$$X' = (\neg c)$$

Example of unit propagation

Given:

- $X = (a \lor \neg b \lor \neg c)$
- $A = \{a \mapsto F, b \mapsto T\}$

We know that:

$$X' = (\neg c)$$

It is obvious that we must assign c false:

$$A' = A \cup \{c \mapsto F\} = \{a \mapsto F, b \mapsto T, c \mapsto F\}$$

$$X'' \mapsto T$$

Evaluating unit propagation

- Unit propagation is an extremely worthwhile inference technique to add to a search procedure
 - Substantial pruning
 - Simple reasoning
- Implementation details have a massive impact on performance, even for a simple and effective technique such as unit propagation

Evaluating inference techniques

- When examining the potential worth of an inference technique, we need to compare:
 - How much of the search space will be pruned (the benefit)
 - How much time and space executing the technique will require (the cost)
- Some of this comparison can be done theoretically
- Implementation details cannot be avoided

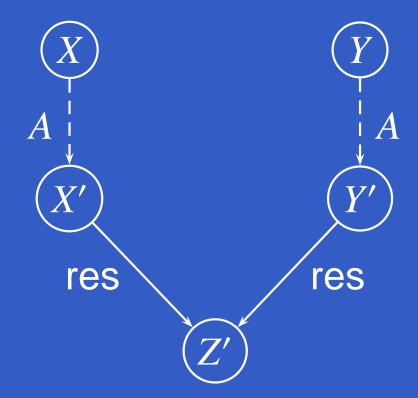
Example of resolution

Example of neighbour resolution

$$X = (\neg a \lor \neg b \lor \neg c)$$

$$Y = (a \lor \neg b \lor \neg c)$$

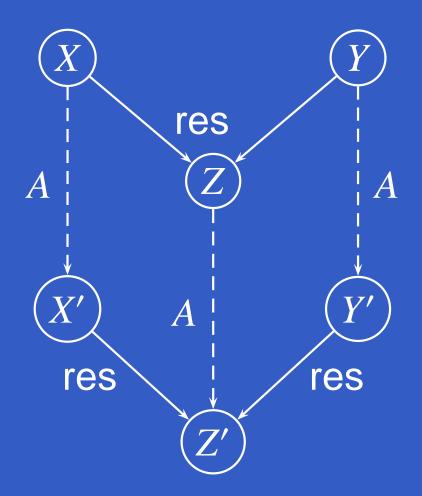
Neighbour resolution



Evaluating neighbour resolution

- Neighbour resolution during search significantly prunes the search space on many problems
- Identifying neighbouring clauses during search takes a great deal of time
- The time cost outweighs the benefit, meaning that this implementation of neighbour resolution during search is not practically beneficial

Preprocessing neighbour resolution



Evaluating preprocessing NR

- Not a completely correct simulation
- Works fairly well for some problem classes, but needs more work

Future work

- More theoretical evaluation
 - How much pruning?
 - How much work compared to unit prop?
 - Which instances respond to a technique?
 - Is it possible to identify an equivalent preprocessing technique?
- Efficiently combining nogood recording with other inference techniques
- Hyper binary resolution
 - Relationship between CSP encodings

Conclusions

- SAT is both theoretically interesting and practically important
- Inference can successfully augment search on SAT problems
- The challenge is to find inference techniques that are cost-effective

Related work

Combining resolution and search:

- Rish and Dechter. Resolution versus search: two strategies for SAT. In SAT2000, IOS Press, 2000.
- van Gelder. Satisfiability testing with more reasoning and less guessing. In Second DIMACS implementation challenge, 1995.
- Cha and Iwama. Adding new clauses for faster local search. In Proc AAAI-96, 1996.

Other worthwhile inference techniques

- Nogood recording
- Conflict-directed backjumping
- Equivalency reasoning
- Restrictions of resolution
- Hyper-resolution
- Variable probing

Future work: improved simulation

In our current implementation of simulated neighbour resolution:

- Subsumption during search is ignored
 - We can mark resolvent clauses and cheaply apply subsumption to just those clauses during search
- Extra resolvents (not corresponding to actual neighbour resolvents) are added
 - We can use knowledge of the branching heuristic to determine which resolvents correspond to actual neighbour resolvents

Future work: improved implementation

- Neighbour resolution during search is slow because identifying neighbouring clauses is expensive
 - We have an improved algorithm for neighbour identification which we plan to implement
- It is not worth applying resolution to some problem classes (e.g. the JNH SATLIB instances)
 - We are developing syntactic methods for identifying some such problem classes

Future work: investigation

- If implied clauses are visible to the branching heuristic, the search tree may actually be grow instead of being pruned
 - We plan to investigate the effect of including the implied clauses, but making the branching heuristic ignore them

My other work

Investigating preprocessing techniques:

- Systematic comparison of existing techniques
- Selecting and evaluating novel techniques
 - E.g. taking first-order techniques and applying them to SAT problems